

Equivalence between operator spreading and information propagation

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Introduction

Operator spreading and the Lieb-Robinson bound

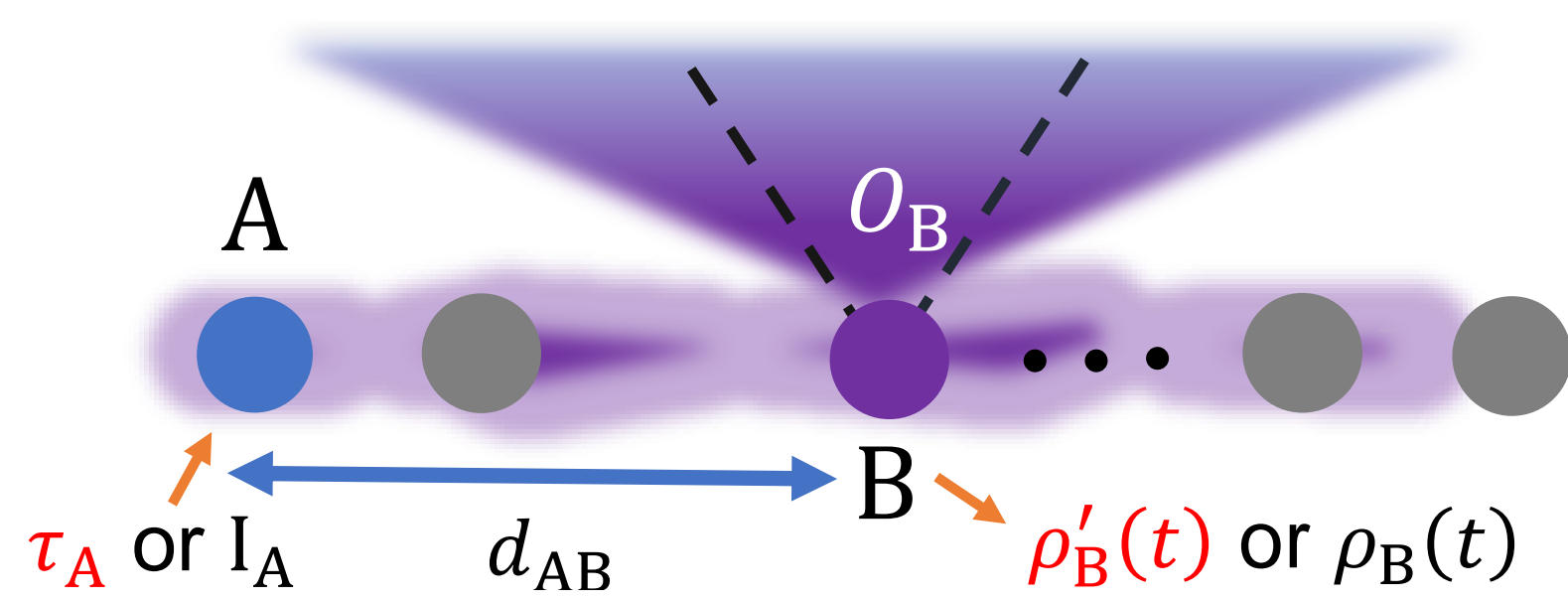


Figure 1: $\varepsilon = \exp[-\text{const}(d_{AB} - v_{LR}t)]$, the Lieb-Robinson bound gives an upper bound on operator spreading, where v_{LR} is called Lieb-Robinson velocity^[1].

Information propagation and the Holevo capacity

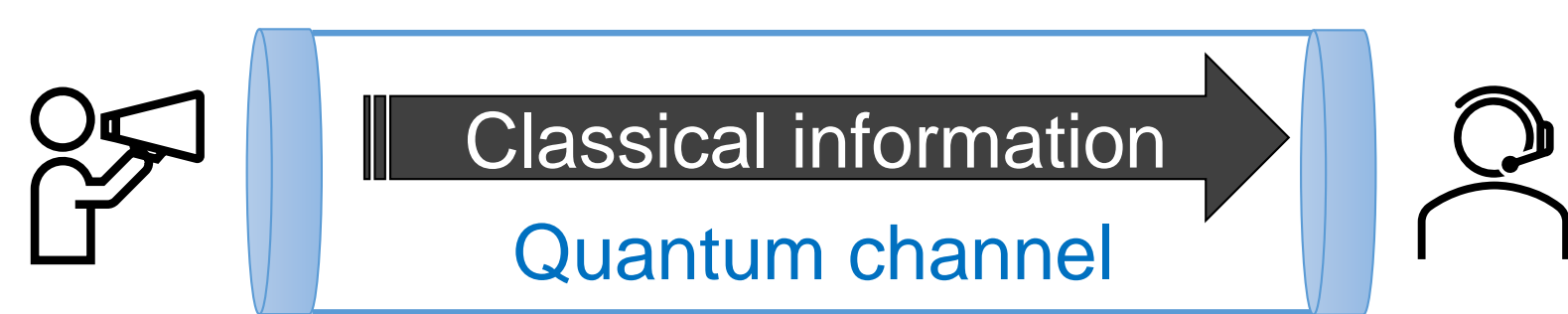


Figure 2: $C_\chi(t)$, the Holevo capacity quantifies the maximum amount of a quantum channel to transmit classical information^[2].

Information propagation as operator spreading

$$C_\chi(t) \propto |\text{Tr}_B \{O_B [\rho_B(t) - \rho'_B(t)]\}| \leq \sup_{\|O_B\|=1} \|O_B(t), \tau_A\|_1 \leq \varepsilon$$

Simple example:

Encoding: $\{\rho_0, I_A\}$ and $\{1 - \rho_0, \tau_A\}$, where I_A and τ_A represent identity and unitary operation to A, respectively.

Unitary evolution: $U_{AB}(t) = \mathcal{T} \exp[-i \int_0^t H_{AB}(\tau) d\tau]$ with any Hamiltonian H_{AB} . We have

$$\rho_B(t) = \text{Tr}_A [U_{AB}(t) I_A \rho_{AB}(0) I_A U_{AB}^\dagger(t)],$$

$$\rho'_B(t) = \text{Tr}_A [U_{AB}(t) \tau_A \rho_{AB}(0) \tau_A^\dagger U_{AB}^\dagger(t)],$$

the reduced density matrices $\rho_B(t)$ and $\rho'_B(t)$ correspond to I_A and τ_A .

Finding tight constraints on the classical capacity of quantum channels

1. $C_\chi \propto T$? The trace distance between stationary states ρ_B and ρ'_B can be written as $T(\rho_B, \rho'_B) := 0.5 \|\rho_B - \rho'_B\|_1$.

2. $C_\chi(t) \propto \varepsilon$? Case of $\rho_B(t)$ and $\rho'_B(t)$ change over time.

3. How to estimate the classical communication rate, $dC_\chi(t)/dt$?

References: [1] E. H. Lieb and D. W. Robinson, 1972. [2] A. Holevo, 1973.

Setup and main results

Our Setup

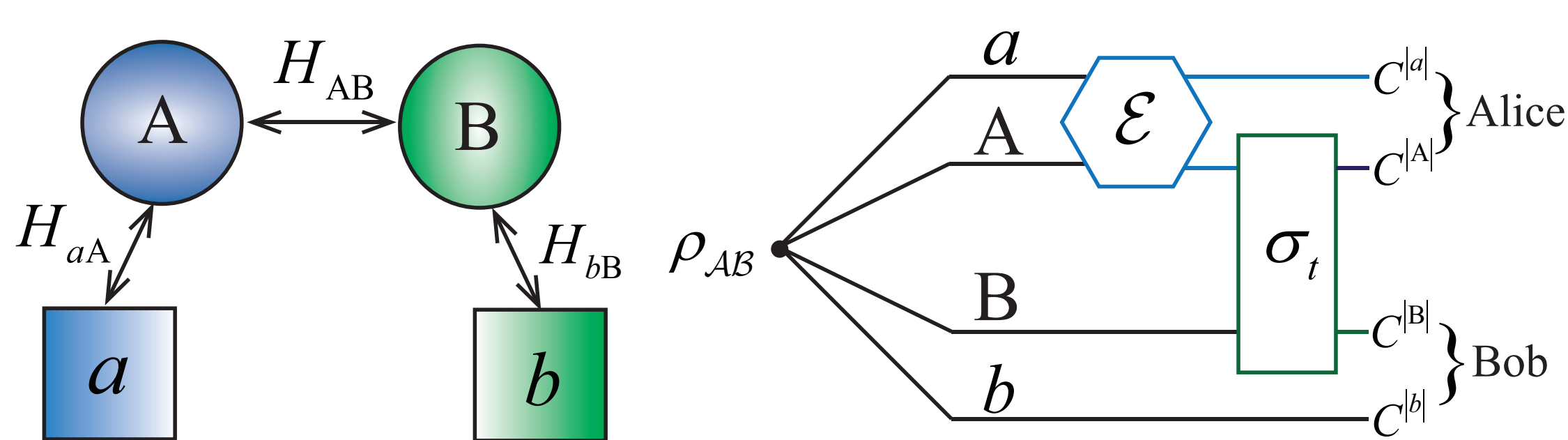


Figure 3: Ancilla-assisted entangling model.

Purification: start from a pure state ρ_{AB} , where $a \cup A = \mathcal{A}$ and $b \cup B = \mathcal{B}$.

Encoding: define \mathcal{E} by a set of CPTP maps acting on \mathcal{A} , $\{p_i, \rho_i = \tau_{A,i}(\rho_{AB})\}$.

A noiseless channel σ_t : $U_{AB}(t) = I_a \otimes \exp(iH_{AB}t) \otimes I_b$.

we have $\rho'_B = \text{tr}_A [\sigma_t \circ \tau_{A,i}(\rho_{AB})] = \text{tr}_A [U_{AB}(t) \tau_{A,i}(\rho_{AB}) U_{AB}^\dagger(t)]$.

Theorem 0. No operator spreading \Leftrightarrow zero information propagation

$$C_\chi(t) = \sum_i p_i S[\rho_B^i(t) || \rho_B(t)] = p_i \rho_B^i(t),$$

for $\forall i, \rho_B^i(t) = \rho_B(t) \Leftrightarrow C_\chi(t) = 0$.

Theorem 1. $C_\chi \propto T(\rho_B^i, \bar{\rho}_B^i)$?

Using Holevo skew divergence and defining the complementary states $\bar{\rho}_B^i$, the time-independent Holevo capacity is bound by trace distance as

$$\sum_i 2p_i(1-p_i)^2 T(\rho_B^i, \bar{\rho}_B^i)^2 \leq C_\chi \leq -\sum_i p_i \log(p_i) T(\rho_B^i, \bar{\rho}_B^i),$$

By applying the improved Pinsker's inequality and defining the measured relative $S_M(\rho_B^i || \bar{\rho}_B^i)$, we obtain $T(\rho_B^i, \bar{\rho}_B^i) \leq \sqrt{1 - \exp[-S_M(\rho_B^i, \bar{\rho}_B^i)]}$.

Theorem 2. $C_\chi(t) \propto \varepsilon$?

Information dynamics means operator spreading, which is limited by the Lieb-Robinson bound. The time-dependent Holevo capacity is bound by LR bound

$$\sum_i 2p_i(1-p_i)^2 \varepsilon^2 \leq C_\chi(t) \leq H(p)\varepsilon$$

where the Shannon entropy $H(p)$ is expressed as $H(p) = -\sum_i p_i \log(p_i)$.

Entanglement capacity

Theorem 3. Entanglement capacity, $dC_\chi(t)/dt$?

We prove that it is captured by the small-incremental-entangling theorem as

$$|dC_\chi(t)/dt| \leq 2\Gamma_t = 2O(1) \|H_{AB}\| \log(d_B), \quad d_B \leq d_A$$

where Γ_t is entangling rate between \mathcal{A} and \mathcal{B} . This result holds for all dimensions d_a, d_b and all states ρ_{AB} of a composite system \mathcal{AB} .

Extension 0-2. α Holevo capacity

By introducing a sandwiched Rényi relative entropy, theorem 0 can be extended to

$$C_\chi^\alpha(t) = \sum_i p_i \mathcal{D}_\alpha[\rho_B^i(t) || \rho_B(t)], \quad \alpha \in (0, +\infty),$$

where

$$\mathcal{D}_\alpha[\rho_B^i(t) || \rho_B(t)] = \frac{1}{1-\alpha} \log \text{Tr}_B \left\{ \left[\rho_B^i(t)^{\frac{1-\alpha}{2\alpha}} \rho_B(t) \rho_B^i(t)^{\frac{1-\alpha}{2\alpha}} \right]^\alpha \right\}.$$

The limiting value of \mathcal{D}_α as $\alpha \rightarrow 1$ back to S . For $C_\chi \propto T(\rho_B^i, \bar{\rho}_B^i)$ and $C_\chi(t) \propto \varepsilon$, we obtain conclusions similar to Theorems 1 and 2.

Extension 3. α entanglement capacity

The generalized entanglement capacity depends on the Rényi order as

$$|dC_\chi^\alpha(t)/dt| \leq \begin{cases} \frac{4\alpha}{|\alpha-1|} \|H_{AB}\| d_B^2, & \alpha < 1, \\ 2O(1) \|H_{AB}\| \log(d_B), & \alpha = 1, \\ \frac{4\alpha}{|\alpha-1|} \|H_{AB}\| d_B^{\frac{2(\alpha-1)}{\alpha}}, & \alpha > 1, \end{cases}$$

Future perspective: Extend our results to noisy channels

References: Cheng Shang, Hayato Kinkawa, and Tomotaka Kuwahara, unpublished 2024.

Acknowledgments

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