# Improving social welfare in non-cooperative games with different types of quantum resources

Alastair A. Abbott<sup>1</sup>, Mehdi Mhalla<sup>2</sup>, and Pierre Pocreau<sup>2</sup>

<sup>1</sup>Univ. Grenoble Alpes, Inria, 38000 Grenoble, France <sup>2</sup>Univ. Grenoble Alpes, CNRS, Grenoble INP, LIG, 38000 Grenoble, France

#### Summary

- Non-cooperative games involving multiple players exhibit equilibria wherein no player has an incentive to deviate from their strategy
- The quality of an equilibrium can be quantified by its social welfare the mean payoff each player receives
- Access to shared quantum resources may allow better cooperation, and hence better equilibria
- We consider two scenarios: in one, players may make measurements directly on a quantum state, while in the other, they delegate the measurement to a referee
- We compare the classes of equilibria obtainable in each setting as well as their maximal social welfare as a function of the bias of the game

#### Non-cooperative games

A non-cooperative game G between n players is defined by:

- $\blacksquare$  A set of **questions**  $T \subseteq \{0,1\}^n$
- lacksquare A prior distribution  $\Pi$  over the questions T
- A set of valid answers  $A = \{0, 1\}^n$
- **A payoff function**  $u_i$  for each player i, with  $u_i(a, t) \in \mathbb{R}$ .
  - We consider payoff functions with the form

$$u_i(a,t) = egin{cases} 0 ext{ if } (a,t) 
otin \mathcal{W} \ v_0 ext{ if } a_i = 0 ext{ and } (a,t) 
otin \mathcal{W} \ v_1 ext{ if } a_i = 1 ext{ and } (a,t) 
otin \mathcal{W}, \end{cases}$$

with  $v_0, v_1 > 0$  and  $\mathcal{W} \subseteq A \times T$  a set of "winning input-output pairs"

Ratio  $v_0/v_1$  controls the bias of game

**Example:** Winning conditions for two 5-player games:  $NC_{00}(C_5)$  and  $NC_{01}(C_5)$  [1]

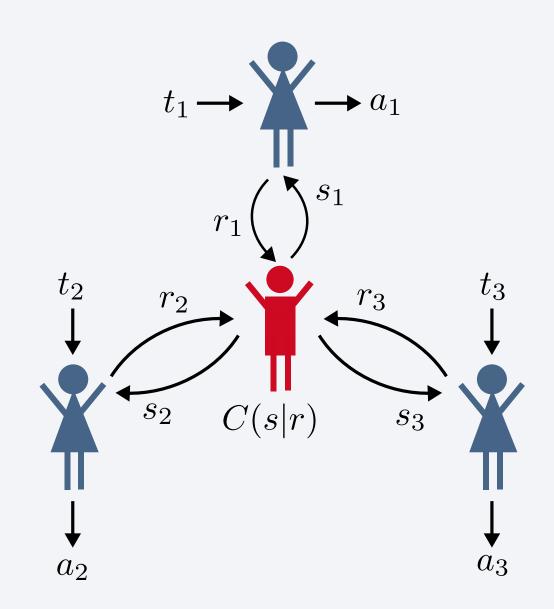
	Winning condition, $NC_{00}(C_5)$	$\begin{array}{c} Question \\ t_1t_1t_2t_3t_5 \end{array}$	Winning condition, $NC_{01}(C_5)$
10000	$a_4 \oplus a_0 \oplus a_1 = 0$	10100	$a_4 \oplus a_0 \oplus a_1 = 0$
01000	$a_0 \oplus a_1 \oplus a_2 = 0$	01010	$a_0 \oplus a_1 \oplus a_2 = 0$
00100	$a_1 \oplus a_2 \oplus a_3 = 0$	00101	$a_1 \oplus a_2 \oplus a_3 = 0$
00010	$a_2 \oplus a_3 \oplus a_4 = 0$	10010	$a_2 \oplus a_3 \oplus a_4 = 0$
00001	$a_3 \oplus a_4 \oplus a_0 = 0$	01001	$a_3 \oplus a_4 \oplus a_0 = 0$
11111	$a_0 \oplus a_1 \oplus a_2 \oplus a_3 \oplus a_4 = 1$	11111	$a_0 \oplus a_1 \oplus a_2 \oplus a_3 \oplus a_4 = 1$

## Solutions and equilibria

- Each player follows a local strategy to produce their answer
- In general, they may also have access to a shared correlation in the form of an advice  $s_i$  provided by a mediator with probability  $C(s_1 \ldots s_n | r_1 \ldots r_n)$
- A solution (set of strategies for each player, defined by functions  $f_i$  and  $g_i$ ) induces a distribution

$$P(a|t) = \sum_{\lambda} \Lambda(\lambda) \sum_{s: \forall i, g_i(t_i, s_i, \lambda_i) = a_i} C(s_1 \dots s_n | f(t_1, \lambda_1) \dots f(t_n, \lambda_n))$$

We can generally consider just deterministic strategies



A solution is a Nash equilibrium if no player can increase their mean payoff by changing their strategy:  $\forall i \ \forall t_i, r_i \in T_i \ \forall \mu_i : T_i \times A_i \to A_i$ ,

$$\sum_{t_{-i}, a_{-i}} u_i(a, t) P(a|t) \Pi(t) \ge \sum_{t_{-i}, a_{-i}} u_i(\mu_i(t_i, a_i) a_{-i}, t) P(a|r_i t_{-i}) \Pi(t)$$

- Nash equilibria play important roles in applications from economics to engineering
- Different correlations C lead to different equilibria: Nash (no correlation), Corr (shared randomness), B.I. (belief invariant, or no-signalling), . . .
- For a game G, the **social welfare** of a solution is

$$SW_G(P) = \frac{1}{n} \sum_{i} \sum_{a,t} u_i(a,t) P(a|t) \Pi(t).$$

### Two types of quantum advices

Question: How can quantum resources lead to new equilibria or improve social welfare?

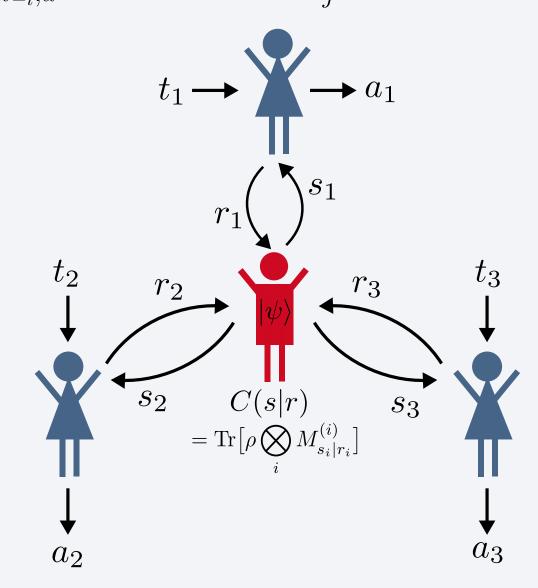
We identify two types of quantum advice leading to different equilibria:

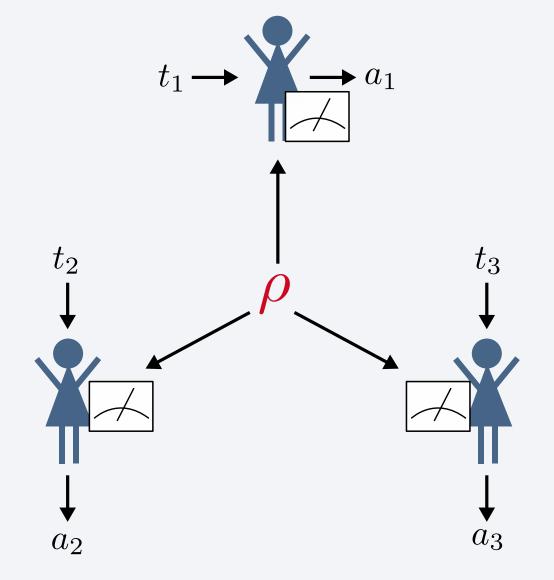
lacktriangle Quantum correlated advice: Advice C is obtained from measurements on a quantum system:

$$C(s|r) = \operatorname{Tr}\left[\rho(M_{s_1|r_1}^{(1)} \otimes \cdots \otimes M_{s_n|r_n}^{(n)})\right]$$

- Measurement delegated to mediator, or performed by parties with quantum "black-boxes"
- Shared quantum state [2]: Each player measures a shared quantum state to determine their output  $a_i$ 
  - Direct access to quantum resource
  - Notion of equilibria modified: a player can deviate by choosing any other local POVM:  $\forall i \, \forall t_i \, \forall N^{(i)} = \{N^{(i)}_{a_i|r_i}\}_{r_i}$

$$\sum_{t_{-i},a} u_i(a,t) \operatorname{Tr} \left( \rho \cdot \bigotimes_j M_{a_j|t_j}^{(j)} \right) \Pi(t) \ge \sum_{t_{-i},a} u_i(a,t) \operatorname{Tr} \left( \rho \cdot \bigotimes_{j \ne i} M_{a_j|t_j}^{(j)} \otimes N_{a_i|t_i}^{(i)} \right) \Pi(t)$$





Classical access –  $Q_{corr}(G)$ 

Quantum access -Q(G)

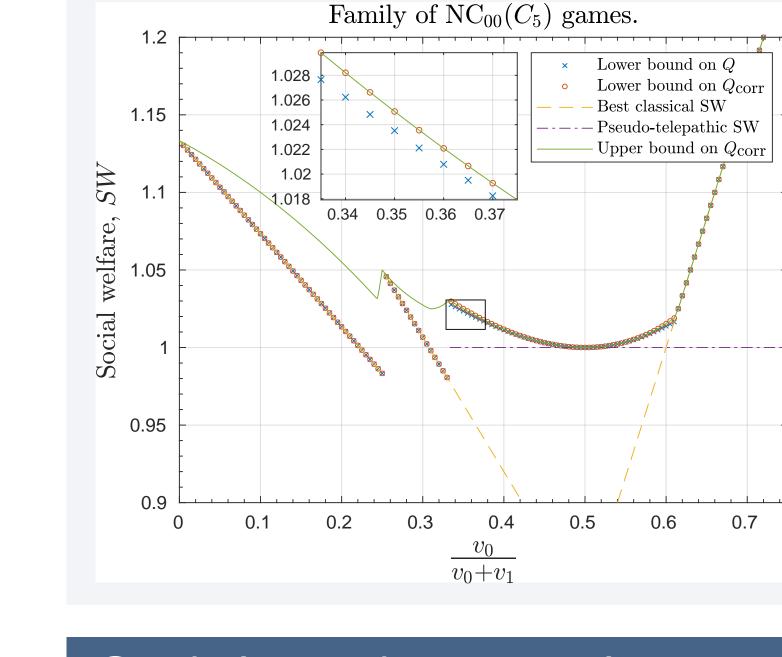
For any game G, the sets of equilibria satisfy

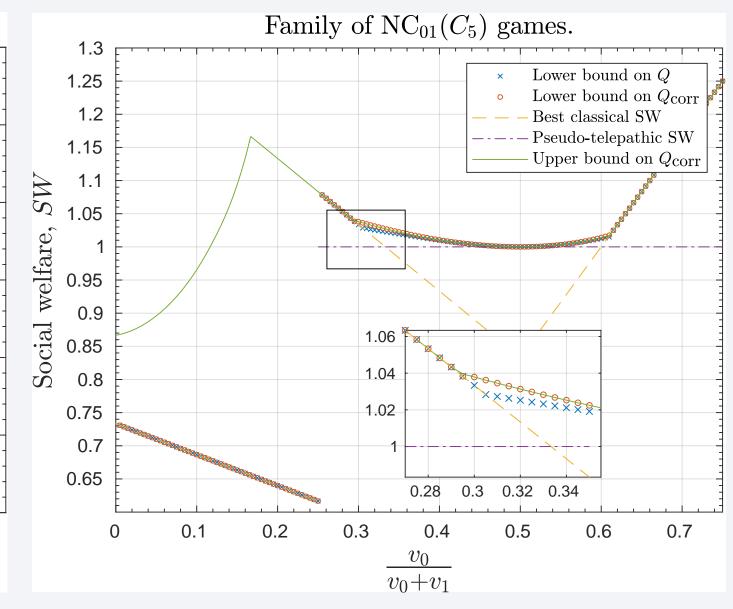
$$\mathsf{Nash}(G) \subset \mathsf{Corr}(G) \subset Q(G) \subset Q_{\mathsf{corr}}(G) \subset \mathsf{B.l.}(G) \subset \mathsf{Comm}(G)$$

- We show using self-testing methods [3] that for some game G,  $Q(G) \subsetneq Q_{corr}(G)$ 
  - Players who delegate quantum measurements can reach more equilibria!

## Results: Social welfare of different solutions

- We optimised the social welfare over different solutions classes for three types of games:  $NC_{00}(C_5)$ ,  $NC_{01}(C_5)$ , and  $NC(C_3)$  (not shown here) [1]
  - Best classical SW: computed exactly
  - **Graph state SW:** pseudo-telepathic equilibria using GHZ states [1]
  - Seesaw lower bound: numerical optimisation by iterating SDPs to find explicit solutions lower-bounding QSW over Q(G) and  $Q_{corr}(G)$
  - NPA upper bound: SDP hierarchy providing dimension-independent upper bound on equilibria in  $Q_{corr}(G)$  [4]





## Conclusions and open questions

- Two different ways to use quantum resources lead to distinct classes of equilibria
- lacksquare A strict separation between Q(G) and  $Q_{\operatorname{corr}}(G)$
- Quantum social welfare can be improved beyond pseudo-telepathic solution
- lacktriangle Method to directly obtain upper bounds on Q(G)?

## References and acknowledgments

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